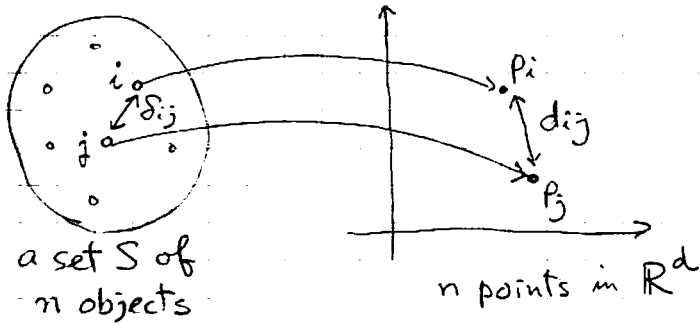


MDS: Multi-Dimensional Scaling



δ_{ij} : dissimilarity between objects i and j

??

d_{ij} : Euclidean distance between two points p_i and p_j

Given an integer $d > 0$ and a matrix $\Delta = [\delta_{ij}]$ representing dissimilarity of every pair of objects, find a mapping of those objects to points in d -dimensional space \mathbb{R}^d so that point-wise distance is approximately equal to the dissimilarity between corresponding objects.

approximation error : $e_{ij} = |\delta_{ij} - d_{ij}|$

objective function to be minimized :

$$f_1 : \sum_{i,j} e_{ij} \rightarrow \min$$

$$f_2 : \sum_{i,j} e_{ij}^2 \rightarrow \min \quad \dots \text{etc.}$$

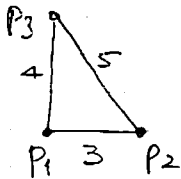
(Example E1)

$$\Delta = \begin{matrix} & \begin{matrix} 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{pmatrix} 0 & 3 & 4 \\ 3 & 0 & 5 \\ 4 & 5 & 0 \end{pmatrix} \end{matrix}$$

$$\delta_{12} = \delta_{21} = 3$$

$$\delta_{13} = \delta_{31} = 4$$

$$\delta_{23} = \delta_{32} = 5$$

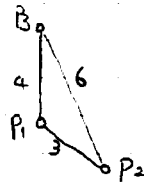


exact embedding in \mathbb{R}^2

(Example E2)

$$\Delta' = \begin{matrix} & \begin{matrix} 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{pmatrix} 0 & 3 & 4 \\ 3 & 0 & 6 \\ 4 & 6 & 0 \end{pmatrix} \end{matrix}$$

Exact embedding
is still possible



(Example E3)

$$\Delta' = \begin{matrix} & \begin{matrix} 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \end{matrix} & \begin{pmatrix} 0 & 3 & 4 \\ 3 & 0 & 8 \\ 4 & 8 & 0 \end{pmatrix} \end{matrix}$$

violates triangular inequality
for $k=i, j, k$

$$d_{ij} + d_{jk} \geq d_{ik}$$

$\left\{ \begin{array}{l} \text{Metric case : every triple satisfies triangular inequality} \\ \text{Non-metric case : otherwise} \end{array} \right.$

Given a set of n objects with associated dissimilarity matrix which is metric (i.e., satisfying triangular inequality), those objects are mapped to points in the $(n-1)$ dimensional space so that $\delta_{ij} = d_{ij}$ for every pair (i, j) .

$x_i = (x_{i1}, x_{i2}, \dots, x_{id})^T$: point in \mathbb{R}^d

\Rightarrow the Euclidian distance d_{ij}

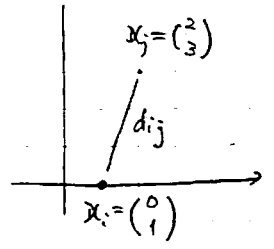
$$d_{ij}^2 = (x_i - x_j)^T (x_i - x_j) \quad (1)$$

(Example E4)

$$x_i = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \quad x_j = \begin{pmatrix} 2 \\ 3 \end{pmatrix}$$

$$\Rightarrow x_i - x_j = \begin{pmatrix} 0-2 \\ 1-3 \end{pmatrix} = \begin{pmatrix} -2 \\ -2 \end{pmatrix}$$

$$d_{ij}^2 = (-2 \ -2) \begin{pmatrix} -2 \\ -2 \end{pmatrix} = (-2)^2 + (-2)^2 = 8$$



Let the inner product matrix be B , where

$$[B]_{ij} = b_{ij} = x_i^T x_j \quad (2)$$

\uparrow Want to find B using known squared distances $\{d_{ij}^2\}$

To find B

Translation does not change pairwise distances

\Downarrow as a natural assumption

$$\sum_{i=1}^n x_{ij} = 0 \quad \text{for } j=1, \dots, d. \quad (3)$$

(means that the centroid of n points is placed at the origin)

$$d_{ij}^2 = (x_i - x_j)^T (x_i - x_j)$$

$$= x_i^T x_i + x_j^T x_j - 2 x_i^T x_j \quad (4)$$

$$\frac{1}{n} \sum_{i=1}^n d_{ij}^2 = \frac{1}{n} \sum_{i=1}^n (x_i^T x_i + x_j^T x_j - 2 x_i^T x_j) \quad (4)$$

$$\text{Here } \sum_{i=1}^n x_i^T x_j = \left(\sum_{i=1}^n x_i^T \right) x_j = 0 \quad (\text{due to (3)})$$

$$\begin{aligned} \therefore \frac{1}{n} \sum_{i=1}^n d_{ij}^2 &= \frac{1}{n} \sum_{i=1}^n x_i^T x_i + \frac{1}{n} \sum_{i=1}^n x_j^T x_j \\ &= \frac{1}{n} \sum_{i=1}^n x_i^T x_i + x_j^T x_j \quad (5) \end{aligned}$$

$$\begin{aligned} \frac{1}{n} \sum_{j=1}^n d_{ij}^2 &= \frac{1}{n} \sum_{j=1}^n (x_i^T x_i + x_j^T x_j - 2 x_i^T x_j) \\ &= x_i^T x_i + \frac{1}{n} \sum_{j=1}^n x_j^T x_j \quad (6) \end{aligned}$$

$$\begin{aligned} \therefore \frac{1}{n^2} \sum_{i=1}^n \sum_{j=1}^n d_{ij}^2 &= \frac{1}{n} \sum_{i=1}^n \left(\frac{1}{n} \sum_{j=1}^n d_{ij}^2 \right) = \frac{1}{n} \sum_{i=1}^n \left(x_i^T x_i + \frac{1}{n} \sum_{j=1}^n x_j^T x_j \right) \\ &= \frac{1}{n} \sum_{i=1}^n x_i^T x_i + \frac{1}{n^2} \sum_{j=1}^n \sum_{i=1}^n x_j^T x_j \\ &= \frac{1}{n} \sum_{i=1}^n x_i^T x_i + \frac{1}{n} \sum_{j=1}^n x_j^T x_j \\ &= \frac{2}{n} \sum_{i=1}^n x_i^T x_i \quad (7) \end{aligned}$$

From (5) and (7),

$$x_j^T x_j = \frac{1}{n} \sum_{i=1}^n d_{ij}^2 - \frac{1}{n} \sum_{i=1}^n x_i^T x_i = \frac{1}{n} \sum_{i=1}^n d_{ij}^2 - \frac{1}{2n^2} \sum_{i=1}^n \sum_{j=1}^n d_{ij}^2$$

Similarly

$$x_i^T x_i = \frac{1}{n} \sum_{j=1}^n d_{ij}^2 - \frac{1}{2n^2} \sum_{i=1}^n \sum_{j=1}^n d_{ij}^2$$

Now, combining (2) and (4), we have

$$\begin{aligned} b_{ij} = x_i^T x_j &= \frac{1}{2} (x_i^T x_i + x_j^T x_j - d_{ij}^2) \\ &= \frac{1}{2} \left(\frac{1}{n} \sum_{j=1}^n d_{ij}^2 - \frac{1}{2n^2} \sum_{i=1}^n \sum_{j=1}^n d_{ij}^2 + \frac{1}{n} \sum_{i=1}^n d_{ij}^2 - \frac{1}{2n^2} \sum_{i=1}^n \sum_{j=1}^n d_{ij}^2 - d_{ij}^2 \right) \\ &= \frac{1}{2} \left(\frac{1}{n} \sum_{i=1}^n d_{ij}^2 + \frac{1}{n} \sum_{j=1}^n d_{ij}^2 - \frac{1}{n^2} \sum_{i=1}^n \sum_{j=1}^n d_{ij}^2 - d_{ij}^2 \right) \end{aligned}$$

$$\therefore b_{ij} = -\frac{1}{2} \left(d_{ij}^2 - \frac{1}{n} \sum_{i=1}^n d_{ij}^2 - \frac{1}{n} \sum_{j=1}^n d_{ij}^2 + \frac{1}{n^2} \sum_{i=1}^n \sum_{j=1}^n d_{ij}^2 \right) \quad (8)$$

Now, define

$$a_{ij} = -\frac{1}{2} d_{ij}^2$$

$$a_{i*} = \frac{1}{n} \sum_{j=1}^n a_{ij} = -\frac{1}{2n} \sum_{j=1}^n d_{ij}^2$$

$$a_{*j} = \frac{1}{n} \sum_{i=1}^n a_{ij} = -\frac{1}{2n} \sum_{i=1}^n d_{ij}^2$$

$$a_{**} = \frac{1}{n^2} \sum_{i=1}^n \sum_{j=1}^n a_{ij} = -\frac{1}{2n^2} \sum_{i=1}^n \sum_{j=1}^n d_{ij}^2 \quad (9)$$

Then, we have

$$b_{ij} = a_{ij} - a_{i*} - a_{*j} + a_{**}. \quad (10)$$

Define matrix A as $[A]_{ij} = a_{ij}$, and hence the inner product matrix B is

$$\boxed{B = HAH} \quad (11)$$

where H is the centering matrix

$$H = I - \frac{1}{n} \mathbf{1}\mathbf{1}^T, \quad \mathbf{1} = \underbrace{(1, 1, \dots, 1)}_n^T$$

e.g.

$$I = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}, \quad \mathbf{1} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}, \quad \mathbf{1}\mathbf{1}^T = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \begin{pmatrix} 1 & 1 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix}$$

$$I - \frac{1}{3} \mathbf{1}\mathbf{1}^T = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} - \frac{1}{3} \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} = \begin{pmatrix} \frac{2}{3} & -\frac{1}{3} & -\frac{1}{3} \\ -\frac{1}{3} & \frac{2}{3} & -\frac{1}{3} \\ -\frac{1}{3} & -\frac{1}{3} & \frac{2}{3} \end{pmatrix}$$

$$HAH = \begin{pmatrix} 1-p & -p & -p \\ -p & 1-p & -p \\ -p & -p & 1-p \end{pmatrix} \begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix} \begin{pmatrix} 1-p & -p & -p \\ -p & 1-p & -p \\ -p & -p & 1-p \end{pmatrix} \quad p = \frac{1}{3}$$

$$= \begin{pmatrix} a_{11} - p \sum_{i=1}^3 a_{i1}, & a_{12} - p \sum_{i=1}^3 a_{i2}, & a_{13} - p \sum_{i=1}^3 a_{i3} \\ \text{---} & \text{---} & \text{---} \\ \text{---} & \text{---} & \text{---} \end{pmatrix} \begin{pmatrix} 1-p & -p & -p \\ -p & 1-p & -p \\ -p & -p & 1-p \end{pmatrix}$$

(1,1) element

$$\begin{aligned}
&= (1-p) \left[a_{11} - p \sum_{i=1}^3 a_{i1} \right] - p \left[a_{12} - p \sum_{i=1}^3 a_{i2} \right] - p \left[a_{13} - p \sum_{i=1}^3 a_{i3} \right] \\
&= a_{11} - p \sum_{i=1}^3 a_{i1} - p \left[a_{11} + a_{12} + a_{13} - p \sum a_{i1} - p \sum a_{i2} - p \sum a_{i3} \right] \\
&= a_{11} - p \sum a_{i1} - p \sum a_{i2} + p^2 \sum \sum a_{ij} \\
&= a_{11} - \frac{1}{n} \sum a_{i1} - \frac{1}{n} \sum a_{i2} + \frac{1}{n^2} \sum \sum a_{ij} \\
&= a_{11} - a_{*1} - a_{1*} + a_{**}
\end{aligned}$$

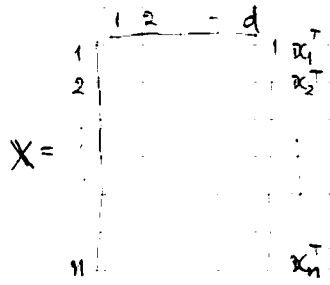
Now, we have

$$[B]_{ij} = x_i^T x_j \quad (2)$$

$$B = HAH^T \quad (11)$$

Setting $X = [x_1, x_2, \dots, x_n]^T$

$$B = XX^T \quad (12)$$



The rank of B, $r(B)$, is then

$$r(B) = r(XX^T) = r(X) = d.$$

Now, B is symmetric, positive semi-definite and of rank d, and hence has d non-negative eigenvalues and $n-d$ zero eigenvalues.

$$Bv_i = \lambda_i v_i, \quad i=1, 2, \dots, n \quad (13)$$

λ_i : eigenvalue

v_i : corresponding eigenvector

Define a matrix V by

$$V = [v_1, v_2, \dots, v_n]$$

Then, from (13) we have

$$BV = V\Lambda, \quad (14)$$

where

$$\Lambda = \begin{pmatrix} \lambda_1 & & 0 \\ & \lambda_2 & \\ 0 & & \ddots \\ & & & \lambda_n \end{pmatrix}$$

with normalizations:

$$v_i^T v_i = 1. \quad (15)$$

And for any $i \neq j$,

$$v_i^T v_j = 0. \quad (16)$$

Therefore, we have

$$VV^T = I.$$

Hence

$$B = BVV^T = V\Lambda V^T$$

$$\text{i.e. } \boxed{B = V\Lambda V^T} \quad (17)$$

Because of the $n-d$ zero eigenvalues, B can be written as

$$B = V_1 \Lambda_1 V_1^T, \quad (18)$$

$$\Lambda_1 = \text{diag}(\lambda_1, \dots, \lambda_d), \quad V_1 = [v_1, \dots, v_d].$$

Hence, as $B = XX^T$, the coordinate matrix X is given by

$$X = V_1 \Lambda_1^{1/2} \tag{19}$$

$$\Lambda_1^{1/2} = \text{diag}(\sqrt{\lambda_1}, \sqrt{\lambda_2}, \dots, \sqrt{\lambda_d})$$

$$X = \begin{matrix} & \begin{matrix} 1 & 2 & \dots & d \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ \vdots \\ n \end{matrix} & \left[\begin{array}{cccc} & & & \\ & & & \\ & & & \\ & & & \end{array} \right] = [v_1, v_2, \dots, v_d] \begin{pmatrix} \sqrt{\lambda_1} & & & 0 \\ & \sqrt{\lambda_2} & & \\ & & \ddots & \\ 0 & & & \sqrt{\lambda_d} \end{pmatrix} \\ & = (\sqrt{\lambda_1} v_1, \sqrt{\lambda_2} v_2, \dots, \sqrt{\lambda_d} v_d)$$

$\sqrt{\lambda_i} v_i$ gives the i -th coordinate values of n points.

(Example E5)

$$\text{dissimilarity matrix} = \begin{pmatrix} 0 & 3 & 4 \\ 3 & 0 & 5 \\ 4 & 5 & 0 \end{pmatrix} = (\delta_{ij})$$

$$a_{ij} = -\frac{1}{2} \delta_{ij}^2 = \begin{pmatrix} 0 & -9/2 & -8 \\ -9/2 & 0 & -25/2 \\ -8 & -25/2 & 0 \end{pmatrix}$$

$$a_{1+} = \frac{1}{3} \sum_{j=1}^3 a_{1j} = \frac{1}{3} (0 - \frac{9}{2} - 8) = -\frac{25}{6} = a_{+1}$$

$$a_{2+} = \frac{1}{3} \sum_{j=1}^3 a_{2j} = \frac{1}{3} (-\frac{9}{2} + 0 - \frac{25}{2}) = -\frac{34}{6} = a_{+2}$$

$$a_{3+} = \frac{1}{3} \sum_{j=1}^3 a_{3j} = \frac{1}{3} (-8 - \frac{25}{2} + 0) = -\frac{41}{6} = a_{+3}$$

$$a_{++} = \frac{1}{3} (a_{1+} + a_{2+} + a_{3+}) = \frac{1}{3} (-\frac{25}{6} - \frac{34}{6} - \frac{41}{6}) = -\frac{100}{18}$$

$$b_{11} = a_{11} - a_{1+} - a_{+1} + a_{++} = 0 + \frac{25}{6} + \frac{25}{6} - \frac{100}{18} = \frac{50}{18} = -\frac{4}{18}$$

$$b_{12} = a_{12} - a_{1+} - a_{+2} + a_{++} = -\frac{9}{2} + \frac{25}{6} + \frac{39}{6} - \frac{100}{18} = \frac{1}{18}(-81 + 75 + 102 - 100)$$

$$b_{13} = a_{13} - a_{1+} - a_{+3} + a_{++} = -8 + \frac{25}{6} + \frac{41}{6} - \frac{100}{18} = \frac{1}{18}(-144 + 75 + 123 - 100) = -\frac{46}{18}$$

$$b_{22} = a_{22} - a_{2+} - a_{+2} + a_{++} = 0 + \frac{36}{6} + \frac{36}{6} - \frac{100}{18} = \frac{1}{18}(102 + 102 - 100) = \frac{104}{18}$$

$$b_{23} = a_{23} - a_{2+} - a_{+3} + a_{++} = -\frac{25}{2} + \frac{39}{6} + \frac{41}{6} - \frac{100}{18} = \frac{1}{18}(-225 + 102 + 123 - 100) = -\frac{100}{18}$$

$$b_{33} = a_{33} - a_{3+} - a_{+3} + a_{++} = 0 + \frac{4}{6} + \frac{41}{6} - \frac{100}{18} = \frac{1}{18}(12 + 123 - 100) = \frac{146}{18}$$

$$B = \frac{1}{18} \begin{pmatrix} 50 & -4 & -46 \\ -4 & 104 & -100 \\ -46 & -100 & 146 \end{pmatrix} = \begin{pmatrix} 2.78, & -0.22, & -2.56 \\ -0.22, & 5.78, & -5.56 \\ -2.56, & -5.56, & 8.11 \end{pmatrix}$$

$$\lambda_1 = 12.96, \lambda_2 = 3.70 \quad \sqrt{\lambda_1} = 3.6, \sqrt{\lambda_2} = 1.92$$

$$v_1^T = (-0.18, -0.59, 0.78)$$

$$v_2^T = (-0.79, 0.56, 0.24)$$

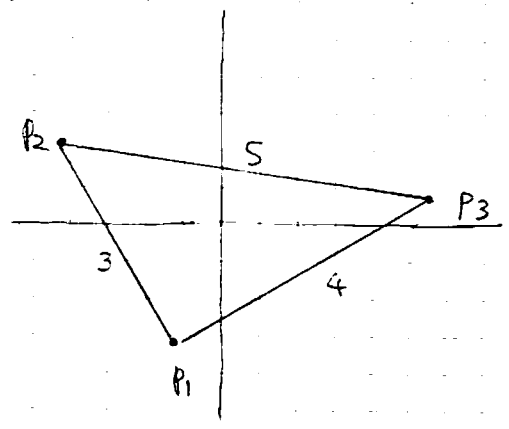
$$\sqrt{\lambda_1} v_1^T = (-0.65, -2.12, 2.81)$$

$$\sqrt{\lambda_2} v_2^T = (-1.52, 1.08, 0.46)$$

$$p_1 = (-0.65, -1.52)$$

$$p_2 = (-2.12, 1.08)$$

$$p_3 = (2.81, 0.46)$$

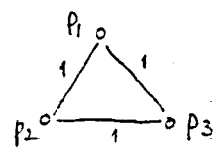


Given a dissimilarity matrix $\{S_{ij}\}$, can we always find an exact embedding if there is?

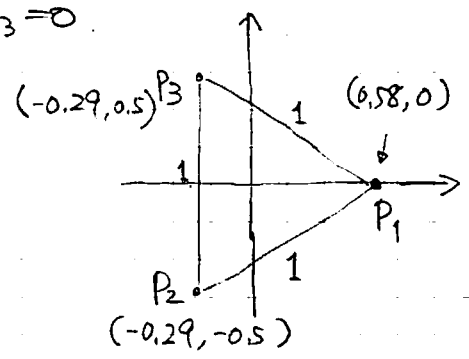
Condition: B is positive semidefinite of rank d
 \Rightarrow exact embedding in \mathbb{R}^d

(Example E6)

$$\Delta = \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix}$$



$$\lambda_1 = \lambda_2 = 0.5, \lambda_3 = 0$$



Question.
 How many dimensions are required?

B has at least one 0 eigenvalue since

$$B\mathbf{1} = HAH\mathbf{1}$$

$$H = \begin{pmatrix} 1 - \frac{1}{n} & -\frac{1}{n} & \dots & -\frac{1}{n} \\ -\frac{1}{n} & 1 - \frac{1}{n} & \dots & -\frac{1}{n} \\ \vdots & \vdots & \ddots & \vdots \\ -\frac{1}{n} & -\frac{1}{n} & \dots & 1 - \frac{1}{n} \end{pmatrix} \quad \mathbf{1} = \begin{pmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{pmatrix}$$

$$H\mathbf{1}_i = (1 - \frac{1}{n}) + \underbrace{(-\frac{1}{n}) + \dots + (-\frac{1}{n})}_{n-1} = 0 \quad \therefore H\mathbf{1} = \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix}$$

$$\therefore B\mathbf{1} = \mathbf{0} = 0\mathbf{1}$$

Suppose we are given a dissimilarity matrix $\{\delta_{ij}\}$ defined by pointwise distances. Then, can we find a configuration such that $d_{ij} = \delta_{ij}$ for every i, j ?

YES by Mardia et al. in 1979.

if B is positive semi-definite of rank d .

$\Rightarrow B = V \Lambda V^T = X X^T$, where

$\Lambda = \text{diag}(\lambda_1, \dots, \lambda_d)$, $X = (x_1, \dots, x_n)^T$, $x_i = \lambda_i^{1/2} v_i$.

$$\begin{aligned}
d_{ij}^2 &= (x_i - x_j)^T (x_i - x_j) = x_i^T x_i + x_j^T x_j - 2 x_i^T x_j \\
&= b_{ii} + b_{jj} - 2 b_{ij} \\
&= (a_{ii} - a_{i*} - a_{*i} + a_{**}) + (b_{jj} - a_{j*} - a_{*j} + a_{**}) \\
&\quad - 2(a_{ij} - a_{i*} - a_{*j} + a_{**}) \\
&= a_{ii} + a_{jj} - 2a_{ij} \quad (\because a_{i*} = a_{*i}, a_{j*} = a_{*j}, \text{etc.}) \\
&= -2a_{ij} \quad (\because a_{ii} = -\frac{1}{2} \delta_{ii}^2 = 0) \\
&= \delta_{ij}^2 \\
\therefore d_{ij}^2 &= \delta_{ij}^2
\end{aligned}$$

Gower (1966) named this method "Principal Coordinates Analysis" (PCO).

in the embedding into \mathbb{R}^{n-1} ,

$d_{ij}^2 = \sum_{k=1}^{n-1} \lambda_k (x_{ik} - x_{jk})^2$ (20)

So, if λ_i is small enough, then their contribution can be neglected.

Principal Coordinates Analysis

- $\{S_{ij}\}$: dissimilarity matrix
- Define a matrix B
- Find d largest eigenvalues $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_d \geq 0$ and their corresponding eigenvectors v_1, \dots, v_d .
- Compute a configuration by

$$X = [\sqrt{\lambda_1} v_1, \sqrt{\lambda_2} v_2, \dots, \sqrt{\lambda_d} v_d]$$

Principal Coordinate Analysis (PCO) is expected to give a good embedding even if there is no exact embedding, but there is no theoretical bound on the performance (e.g., as an approximation ratio).

Metric Least Squares Scaling

$$\text{loss function } S = \sum_{i < j} \frac{1}{\delta_{ij}} (d_{ij} - \delta_{ij})^2 / \sum_{i < j} \delta_{ij} \quad (21)$$

↓

to be minimized.

Now, $d_{ij}^2 = \sum_{k=1}^d (x_{ik} - x_{jk})^2$, and hence

$$\begin{aligned} \frac{\partial d_{ij}}{\partial x_{tk}} &= \frac{\partial}{\partial x_{tk}} \sqrt{(x_{i1} - x_{j1})^2 + \dots + (x_{id} - x_{jd})^2} \\ &= \frac{1}{2} \cdot \frac{1}{d_{ij}} \cdot \left[2(x_{ik} - x_{jk}) \delta^{it} - 2(x_{ik} - x_{jk}) \delta^{jt} \right] \\ &= \frac{1}{d_{ij}} (x_{ik} - x_{jk}) (\delta^{it} - \delta^{jt}) \end{aligned} \quad (22)$$

$$\text{where } \delta^{rs} = \begin{cases} 1 & \text{if } r=s \\ 0 & \text{if } r \neq s \end{cases}$$

$$\begin{aligned} \frac{\partial S}{\partial x_{tk}} &= \frac{\partial}{\partial x_{tk}} \left[\sum_{i < j} \frac{1}{\delta_{ij}} (d_{ij} - \delta_{ij})^2 / \sum_{i < j} \delta_{ij} \right] \\ &= \left[\sum_{i < j} \frac{1}{\delta_{ij}} 2(d_{ij} - \delta_{ij}) \frac{\partial d_{ij}}{\partial x_{tk}} \right] / \sum_{i < j} \delta_{ij} \\ &= \left[\sum_{i < j} \frac{1}{\delta_{ij}} 2(d_{ij} - \delta_{ij}) \frac{1}{d_{ij}} (x_{ik} - x_{jk}) (\delta^{it} - \delta^{jt}) \right] / \sum_{i < j} \delta_{ij} \\ &= \left(\frac{2}{\sum_{i < j} \delta_{ij}} \right) \frac{\partial}{\partial x_{tk}} \left(\sum_{i=1}^n \frac{1}{\delta_{it}} (d_{it} - \delta_{it})^2 \right) \\ &= \left(\frac{2}{\sum_{i < j} \delta_{ij}} \right) \sum_{i=1}^n \frac{d_{it} - \delta_{it}}{\delta_{it}} \cdot \frac{1}{d_{it}} (x_{ik} - x_{tk}) (\delta^{it} - \delta^{tk}) \\ &= \left(\frac{2}{\sum_{i < j} \delta_{ij}} \right) \sum_{i=1}^n \frac{d_{it} - \delta_{it}}{\delta_{it} d_{it}} (x_{tk} - x_{ik}). \end{aligned}$$

The equations

$$\frac{\partial S}{\partial \chi_{tk}} = \left(\frac{2}{\sum_{i=1}^n \delta_{ij}} \right) \sum_{i=1}^n \frac{d_{it} - \delta_{it}}{d_{it} \delta_{it}} (\chi_{tk} - \chi_{ik}) = 0 \quad (23)$$

$$k = 1, \dots, n, \quad t = 1, \dots, d$$

have to be solved numerically.

Sammon (1969) uses a steepest descent method, so that if $\chi_{tk}^{(m)}$ is the m -th iteration in minimizing S , then

$$\chi_{tk}^{(m+1)} = \chi_{tk}^{(m)} - MF \frac{\partial S}{\partial \chi_{tk}} / \left| \frac{\partial^2 S}{\partial \chi_{tk}^2} \right|, \quad (24)$$

where MF is Sammon's magic factor to optimize convergence and was chosen as 0.3 or 0.4.

Least absolute residuals.

$$LAR = \sum_{i=1}^n w_{ij} |d_{ij} - \delta_{ij}|. \quad (25)$$

where w_{ij} are weights.

Unidimensional Scaling

in one dimension

$$d_{ij} = |x_i - x_j| \quad (26)$$

and the loss function to be minimized is

$$S = \sum_{i < j} (\delta_{ij} - |x_i - x_j|)^2 \quad (27)$$

x is a local minimum of S if and only if

$$x_i = \frac{1}{n} \sum_{j=1}^n \delta_{ij} \operatorname{sign}(x_i - x_j), \quad i=1, 2, \dots, n. \quad (28)$$

where
$$\operatorname{sign}(a-b) = \begin{cases} 1 & a > b \\ 0 & a = b \\ -1 & a < b \end{cases} \quad (29)$$

Special case : ordering is fixed $x_1 \geq x_2 \geq \dots \geq x_n$.

$$S = \sum_{i < j} (\delta_{ij} - (x_i - x_j))^2 \quad (30)$$

$$\begin{aligned} \frac{\partial S}{\partial x_k} &= \sum_{k=i < j} (-2\delta_{kj} + 2(x_k - x_j)) + \sum_{i < k=j} (2\delta_{ik} + 2(x_k - x_i)) \\ &= \sum_{j=k+1}^n (-2\delta_{kj} + 2(x_k - x_j)) + \sum_{i=1}^{k-1} (2\delta_{ik} + 2(x_k - x_i)) \\ &= 2(n-1)x_k - \sum_{i \neq k} 2x_i + \sum_{i=1}^{k-1} 2\delta_{ik} - \sum_{i=k+1}^n 2\delta_{ik} \end{aligned}$$

So, the equations to be satisfied are

$$(n-1)x_k - \sum_{i \neq k} x_i = - \sum_{i=1}^{k-1} \delta_{ik} + \sum_{i=k+1}^n \delta_{ik} \quad (32)$$

Since we have at most $(n-1)$ independent equations, the system of equations is expressed in the following matrix form:

$$\begin{pmatrix} n-1 & -1 & \dots & -1 \\ -1 & n-1 & & \vdots \\ \vdots & -1 & \ddots & \vdots \\ -1 & -1 & & n-1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_{n-1} \end{pmatrix} = \begin{pmatrix} C_1 \\ C_2 \\ \vdots \\ C_{n-1} \end{pmatrix}$$

where $C_i = -\sum_{j=1}^{i-1} \delta_{ij} + \sum_{j=i+1}^n \delta_{ij}$, $i=1, \dots, n$.

Fortunately, the matrix has an inverse matrix. In fact,

$$\begin{pmatrix} \frac{2}{n} & \frac{1}{n} & \dots & \frac{1}{n} \\ \frac{1}{n} & \frac{2}{n} & & \vdots \\ \vdots & \vdots & \ddots & \vdots \\ \frac{1}{n} & \frac{1}{n} & & \frac{2}{n} \end{pmatrix} \begin{pmatrix} n-1 & -1 & \dots & -1 \\ -1 & n-1 & & \vdots \\ \vdots & -1 & \ddots & \vdots \\ -1 & -1 & \dots & n-1 \end{pmatrix} = \begin{pmatrix} 1 & & & 0 \\ & 1 & & 0 \\ & & \ddots & \\ 0 & & & 1 \end{pmatrix}$$

\therefore diagonal = $\frac{2}{n} \times (n-1) + (n-2) \times \frac{1}{n} \times (-1) = \frac{1}{n} (2n-2-n+2) = 1$.

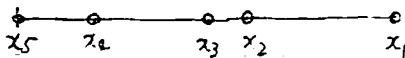
non-diagonal = $\frac{2}{n} \times (-1) + \frac{1}{n} \times (n-1) + (n-3) \times \frac{1}{n} \times (-1)$
 $= \frac{1}{n} (-2+n-1-n+3) = 0$

Thus, we have

$$\begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_{n-1} \end{pmatrix} = \begin{pmatrix} \frac{2}{n} & \frac{1}{n} & \dots & \frac{1}{n} \\ \frac{1}{n} & \frac{2}{n} & & \vdots \\ \vdots & \vdots & \ddots & \vdots \\ \frac{1}{n} & \frac{1}{n} & & \frac{2}{n} \end{pmatrix} \begin{pmatrix} C_1 \\ C_2 \\ \vdots \\ C_{n-1} \end{pmatrix}$$

$$x_n = \frac{1}{n-1} \left[\sum_{i=1}^{n-1} x_i - \sum_{i=1}^{n-1} \delta_{in} \right]$$

(Example E7)



$$(\delta_{ij}) = \begin{pmatrix} 1 & 3 & 4 & 5 \\ 0 & 4 & 5 & 8 & 10 \\ 7 & 0 & 1 & 4 & 6 \\ 5 & 1 & 0 & 3 & 5 \\ 8 & 4 & 3 & 0 & 2 \\ 10 & 6 & 5 & 2 & 0 \end{pmatrix}$$

$$C_1 = -\sum_{j=1}^0 \delta_{1j} + \sum_{j=2}^5 \delta_{1j} = 27$$

$$C_2 = -\sum_{j=1}^1 \delta_{2j} + \sum_{j=3}^5 \delta_{2j} = 4 + 11 = 7$$

$$C_3 = -\sum_{j=1}^2 \delta_{3j} + \sum_{j=4}^5 \delta_{3j} = -6 + 8 = 2$$

$$C_4 = -\sum_{j=1}^3 \delta_{4j} + \sum_{j=5}^5 \delta_{4j} = -15 + 2 = -13$$

$$C_5 = -\sum_{j=1}^4 \delta_{5j} = -23$$

$$C_1 + C_2 + C_3 + C_4 + C_5 = 27 + 7 + 2 - 13 - 23 = 36 - 36 = 0$$

$$x_1 = \frac{1}{n} C_1 + \frac{1}{n} \sum C_i = \frac{27}{5}$$

$$x_2 = \frac{1}{n} C_2 + \frac{1}{n} \sum C_i = \frac{7}{5}$$

$$x_3 = \frac{1}{n} C_3 + \frac{1}{n} \sum C_i = \frac{2}{5}$$

$$x_4 = \frac{1}{n} C_4 + \frac{1}{n} \sum C_i = -\frac{13}{5}$$

$$x_5 = \frac{1}{4} [x_1 + x_2 + x_3 + x_4 + C_5] = \frac{1}{4} \left[\frac{1}{5} (27 + 7 + 2 - 13) - 23 \right] = \frac{1}{4} \left(\frac{23}{5} - 23 \right) = -\frac{23}{5}$$

(Example E8)

$$\delta_{ij} = \begin{cases} 1 & \text{if } i \neq j \\ 0 & \text{if } i = j \end{cases}$$

$$C_1 = n - 2$$

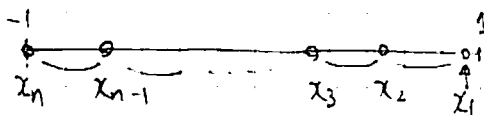
$$C_2 = n - 2 - 2 = n - 4$$

$$C_i = n - 2 - 2*(i-1) = n - 2i$$

$$C_n = n - 2 - 2(n-1) = -n$$

$$\sum_{i=1}^n C_i = \sum_{i=1}^n (n - 2i) = n(n-1) - 2 \cdot \frac{1}{2} n(n-1) = 0$$

$$\therefore x_i = \frac{1}{n} C_i = \frac{1}{n} (n - 2i) = 1 - \frac{2i}{n}, \quad x_1 = 1 - \frac{2}{n}, \quad x_n = 1 - 2 = -1$$



$$\begin{aligned} b_{ij} &= a_{ij} - a_{i*} - a_{*j} + a_{**} \\ &= -\frac{1}{2} + \frac{1}{2n}(n-1) + \frac{1}{2n}(n-1) - \frac{1}{2n}(n-1) = -\frac{1}{2n} \quad \text{if } i \neq j \end{aligned}$$

$$\begin{aligned} b_{ii} &= a_{ii} - a_{i*} - a_{*i} + a_{**} \\ &= 0 + \frac{1}{2n}(n-1) + \frac{1}{2n}(n-1) - \frac{1}{2n}(n-1) = \frac{1}{2} - \frac{1}{2n} \end{aligned}$$

If we set $\varepsilon = -\frac{1}{2n}$, we have

$$B = \begin{pmatrix} \frac{1}{2} + \varepsilon & & & & \\ & \frac{1}{2} + \varepsilon & & & \\ & & \ddots & & \\ \varepsilon & & & \ddots & \\ & & & & \frac{1}{2} + \varepsilon \end{pmatrix}$$

• The matrix B has an eigenvalue $\frac{1}{2}$:

for any vector $v = (v_1, \dots, v_n)^T$ such that $\sum_{i=1}^n v_i = 0$,

$$Bv = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}, \quad x_i = \frac{1}{2}v_i + \varepsilon \sum_{j=1}^n v_j = \frac{1}{2}v_i$$

$$\therefore Bv = \frac{1}{2}v \quad \Rightarrow \lambda = \frac{1}{2} \text{ is an eigenvalue}$$

• The matrix B has an eigenvalue 0 :

for the vector $\mathbf{1} = (1, \dots, 1)^T$,

$$B\mathbf{1} = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}, \quad x_i = \frac{1}{2} + \varepsilon n = \frac{1}{2} + n \times \left(-\frac{1}{2n}\right) = 0.$$

$$\therefore B\mathbf{1} = 0\mathbf{1} \quad \Rightarrow \lambda = 0 \text{ is an eigenvalue}$$

• B has no other eigenvalues.

Let λ and $v = (v_1, \dots, v_n)^T$ be a pair of eigenvalue and eigen vector of the matrix B .

Then, we have

$$Bv = \lambda v.$$

$$\Rightarrow \lambda v_i = \frac{1}{2}v_i + \varepsilon \sum_{j=1}^n v_j, \quad i=1, 2, \dots, n, \quad \varepsilon = -\frac{1}{2n}.$$

Case 1: $\sum_{j=1}^n v_j = 0$

$$\lambda v_i = \frac{1}{2}v_i, \quad \therefore \lambda = \frac{1}{2} \text{ (eigenvalue)}$$

Case 2: $\sum_{j=1}^n v_j \neq 0$

we must have

$$\lambda v_i = \frac{1}{2}v_i - \frac{1}{2n} \sum_{j=1}^n v_j.$$

For any $v_i \neq 0$, we must have

$$\lambda = \frac{1}{2} - \frac{1}{2n v_i} \sum_{j=1}^n v_j,$$

independently of i . Therefore, we must have

$v_1 = v_2 = \dots = v_n$. Then, we have

$$\lambda = \frac{1}{2} - \frac{1}{2n v_i} \cdot n v_i = \frac{1}{2} - \frac{1}{2} = 0. \text{ (eigenvalue)}$$

In this special case, eigenvector corresponding to the eigenvalue $\lambda_1 = 1/2$ is not uniquely determined.

Any vector $v = (v_1, \dots, v_n)^T$ with $\sum_{j=1}^n v_j = 0$ can be an eigenvector of $\lambda = 1/2$.

• Special case : $n = 3$

Let three points be $p_1(x_1, y_1)$, $p_2(x_2, y_2)$, and $p_3(x_3, y_3)$.

$$\lambda_1 = \lambda_2 = 1/2, \quad v_1 = (x_1, x_2)^T \text{ and } v_2 = (y_1, y_2)^T$$

condition: $x_1 + x_2 + x_3 = 0$ and $y_1 + y_2 + y_3 = 0$.

$$B = \begin{pmatrix} \frac{1}{2} + \epsilon & \epsilon \\ \epsilon & \frac{1}{2} + \epsilon \end{pmatrix} = \begin{pmatrix} \frac{1}{2} - \frac{1}{6} & -\frac{1}{6} \\ -\frac{1}{6} & \frac{1}{2} - \frac{1}{6} \end{pmatrix} = \begin{pmatrix} \frac{1}{3} & -\frac{1}{6} \\ -\frac{1}{6} & \frac{1}{3} \end{pmatrix}$$

$$B u_1 = \begin{pmatrix} \frac{1}{3} x_1 - \frac{1}{6} x_2 \\ -\frac{1}{6} x_1 + \frac{1}{3} x_2 \end{pmatrix} = \frac{1}{2} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix}$$

$$\therefore \frac{1}{3} x_1 - \frac{1}{6} x_2 = \frac{1}{2} x_1 \rightarrow x_1 = -x_2$$

$$-\frac{1}{6} x_1 + \frac{1}{3} x_2 = \frac{1}{2} x_2 \rightarrow x_1 = -x_2$$

Similarly, from $B u_2 = \frac{1}{2} u_2$, we have

$$y_1 = -y_2$$

Thus, we have

$$\textcircled{1} x_1 + x_2 + x_3 = 0$$

$$\textcircled{2} x_1 = -x_2 \Rightarrow x_1 = -x_2 \text{ and } x_3 = 0$$

$$\textcircled{3} y_1 + y_2 + y_3 = 0 \Rightarrow y_1 = -y_2 \text{ and } y_3 = 0$$

$$\textcircled{4} y_1 = -y_2$$

• Two vectors must be orthogonal

$$w_1^T u_2 = (x_1 \ x_2 \ x_3) \begin{pmatrix} y_1 \\ y_2 \\ y_3 \end{pmatrix} = x_1 y_1 + x_2 y_2 + x_3 y_3 = 1$$

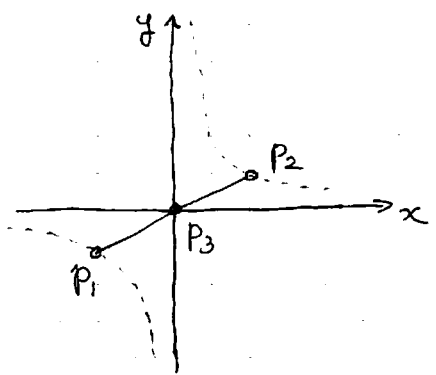
$$\Rightarrow x_1 y_1 + (-x_1)(-y_1) + 0 \cdot 0 = 2x_1 y_1 = 1$$

• Summarize the constraints, we have

$$\textcircled{1} x_1 = -x_2, x_3 = 0$$

$$\textcircled{2} y_1 = -y_2, y_3 = 0$$

$$\textcircled{3} x_1 y_1 = \frac{1}{2}$$



Various Objective Functions

$$S_1 = \sum_{i < j} |d_{ij}^2 - \delta_{ij}^2|$$

$$S_2 = \sum_{i < j} (d_{ij} - \delta_{ij})^2$$

$$S_3 = \sum_{i < j} |d_{ij} - \delta_{ij}|$$

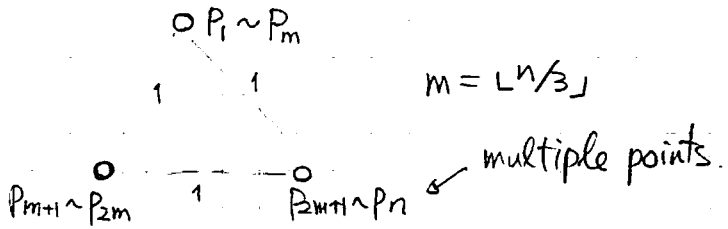
$$S_4 = \sum_{i < j} (d_{ij}^2 - \delta_{ij}^2)^2$$

many other variations, e.g.,

$$S'_2 = \sum_{i < j} \frac{1}{\delta_{ij}} (d_{ij} - \delta_{ij})^2 / \sum_{i < j} \delta_{ij}$$

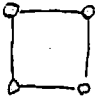
$$S_1 = \sum_{i < j} |d_{ij}^2 - \delta_{ij}^2|$$

Special case : $\delta_{ij} = \begin{cases} 1 & \text{if } i \neq j \\ 0 & \text{if } i = j \end{cases}$.

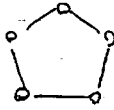


$$S_2 = \sum_{i < j} (d_{ij} - \delta_{ij})^2$$

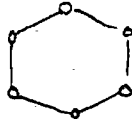
$$\delta_{ij} = \begin{cases} 1 & \text{if } i \neq j \\ 0 & \text{if } i = j \end{cases}$$



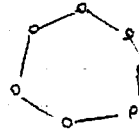
n=4



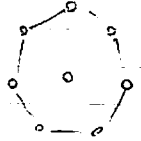
n=5



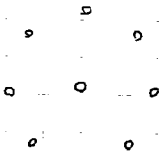
n=6



n=7



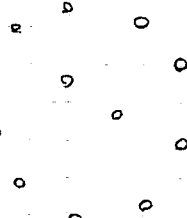
n=8



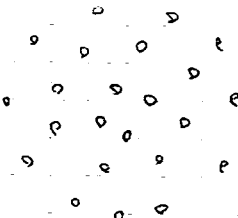
n=9



n=11



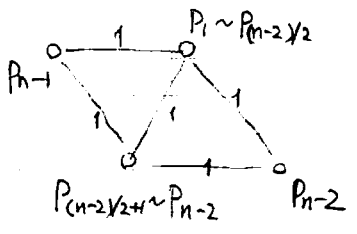
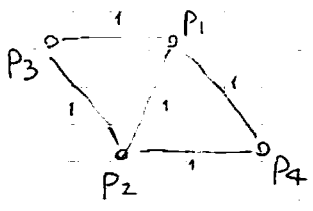
n=12



n=13

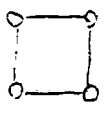
$$S_3 = \sum_{i < j} |d_{ij} - \delta_{ij}|$$

$$\delta_{ij} = \begin{cases} 1 & \text{if } i \neq j \\ 0 & \text{if } i = j \end{cases}$$

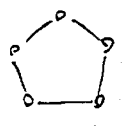


$$S_4 = \sum_{i < j} (d_{ij}^2 - \delta_{ij}^2)^2$$

$$\delta_{ij} = \begin{cases} 1 & \text{if } i \neq j \\ 0 & \text{if } i = j \end{cases}$$



n=4



n=5

...



n=10

Inserting one point in the midst of many points

Problem: Given n points p_1, \dots, p_n in the plane and n real values s_1, \dots, s_n , find a point p that minimizes

$$\sum_{i=1}^n |d_i^2 - s_i^2|,$$

where d_i is the distance from p to p_i .

Note!

A similar problem with an objective function

$$\sum_{i=1}^n |d_i - s_i|$$

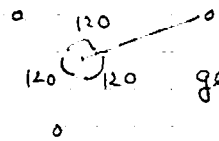
is hard.

∴ Consider a special case where $s_i = 0, i = 1, \dots, n$.

This problem is to find a point that minimizes the sum of distances to the existing points.

→ It is called "Geometric Median"

For three points



geometric median for three points

↳ known as "Fermat point"

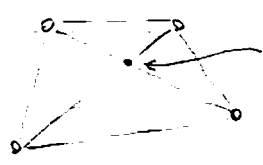
This problem is called the Fermat-Weber problem.

For four points



→ geometric center coincides with the internal point.

Case 1 : one point is inside the triangle formed by the remaining three points



intersection of two diagonals → geometric center

Case 2 : convex positions

For n points.

No formula is known to calculate the geometric median.

Good news : the sum of distances is a convex function



local optimum = global optimum.

Problem is hard even for 5 points.

Approximation algorithm (Bose-Makeshwari-Morin'03).

ϵ -approximation algorithms

$O(n \log n)$ time : deterministic

$O(n)$ time : randomized

FACT: The search space is convex.

Proof: the distance to each existing point p_i is convex.
($f_i = d(p, p_i)$)

- sum of convex functions is convex
- the sum of distances to the existing points is convex. $\sum d(p, p_i)$

Problem: Given n points p_1, \dots, p_n and n real values $\delta_1, \dots, \delta_n$, find a point p that minimizes

$$\sum_{i=1}^n |d_i^2 - \delta_i^2|,$$

where d_i is the distance from p to p_i .

• Special case $\delta_i = 0, i=1, \dots, n$.

$$\text{Find a point } p : \sum_{i=1}^n d_i^2 \rightarrow \min$$

Again, the search space is convex.

$$S = \sum_{i=1}^n d_i^2 = \sum_{i=1}^n (x-x_i)^2 + (y-y_i)^2$$

At a global optimum, we must have

$$\frac{\partial S}{\partial x} = \frac{\partial S}{\partial y} = 0.$$

Thus, we have

$$\frac{\partial S}{\partial x} = 2 \sum_{i=1}^n (x-x_i) = 0 \quad \therefore x = \frac{1}{n} \sum x_i$$

$$\frac{\partial S}{\partial y} = 2 \sum_{i=1}^n (y-y_i) = 0 \quad \therefore y = \frac{1}{n} \sum y_i$$

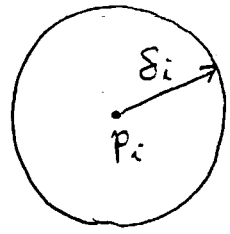
So, the centroid is the global optimum.

Original problem: $S = \sum_{i=1}^n |d_i^2 - \delta_i^2| \rightarrow \min$

unfortunately, the function $|d_i^2 - \delta_i^2|$ is not convex.
So, the search space is not convex.

How to calculate $|d_i^2 - \delta_i^2|$

Draw a circle C_i of radius δ_i and center at p_i



In the interior of C_i

$$|d_i^2 - \delta_i^2| = \delta_i^2 - d_i^2$$

In the exterior of C_i

$$|d_i^2 - \delta_i^2| = d_i^2 - \delta_i^2$$

For n points

Draw n circles C_1, \dots, C_n

\rightarrow arrangement of those circles
 $O(n^2)$ cells.

At each cell R_i we have a quadratic polynomial

$$a_i(x^2 + y^2) + b_i x + c_i y + d_i$$

$$\frac{\partial S}{\partial x} = 2a_i x + b_i = 0 \quad x = -\frac{b_i}{2a_i}$$

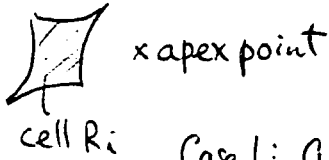
$$\frac{\partial S}{\partial y} = 2a_i y + c_i = 0 \quad y = -\frac{c_i}{2a_i}$$

$(-\frac{b_i}{2a_i}, -\frac{c_i}{2a_i})$: apex point.

If the apex point is within the cell, the apex point is a candidate of an optimal point.

If the apex point is outside the cell

(31)



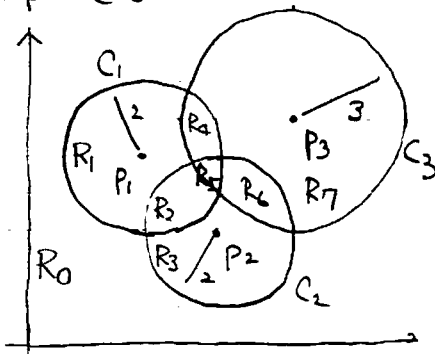
Case 1: $A_i > 0$

find a point on the boundary of R_i that is nearest to the apex point.
→ candidate for the cell.

Case 2: $A_i < 0$

find a point on the boundary of R_i that is farthest to the apex point.

(Example E8)



$$P_1(3,5) \quad \delta_1 = 2$$

$$P_2(5,3) \quad \delta_2 = 2$$

$$P_3(7,6) \quad \delta_3 = 3$$

Cells are coded by three signs $(\sigma_1, \sigma_2, \sigma_3)$

$$\sigma_i = \begin{cases} 1 & \text{if the cell is outside the circle } C_i \\ -1 & \text{otherwise} \end{cases}$$

Then, the objective function

$$\begin{aligned} S &= \sigma_1 [(x-3)^2 + (y-5)^2 - 2^2] + \sigma_2 [(x-5)^2 + (y-3)^2 - 2^2] + \sigma_3 [(x-7)^2 + (y-6)^2 - 3^2] \\ &= (\sigma_1 + \sigma_2 + \sigma_3)(x^2 + y^2) - 2(3\sigma_1 + 5\sigma_2 + 7\sigma_3)x - 2(5\sigma_1 + 3\sigma_2 + 6\sigma_3)y \\ &\quad + \text{const.} \end{aligned}$$

$$\frac{\partial S}{\partial x} = 2(\sigma_1 + \sigma_2 + \sigma_3)x - 2(3\sigma_1 + 5\sigma_2 + 7\sigma_3) = 0$$

$$\frac{\partial S}{\partial y} = 2(\sigma_1 + \sigma_2 + \sigma_3)y - 2(5\sigma_1 + 3\sigma_2 + 6\sigma_3) = 0$$

$$\therefore (x, y) = \left(\frac{3\sigma_1 + 5\sigma_2 + 7\sigma_3}{\sigma_1 + \sigma_2 + \sigma_3}, \frac{5\sigma_1 + 3\sigma_2 + 6\sigma_3}{\sigma_1 + \sigma_2 + \sigma_3} \right)$$

cell $R_0(1, 1, 1)$ $x = 15/3 = 5$ $y = 14/3$ $(x, y) = (5, 14/3)$

cell $R_1(-1, 1, 1)$ $x = 9/1 = 9$ $y = 4/1 = 4$ $(x, y) = (9, 4)$

cell $R_2(-1, -1, 1)$ $x = -1/(-1) = 1$ $y = -2/(-1) = 2$ $(x, y) = (1, 2)$

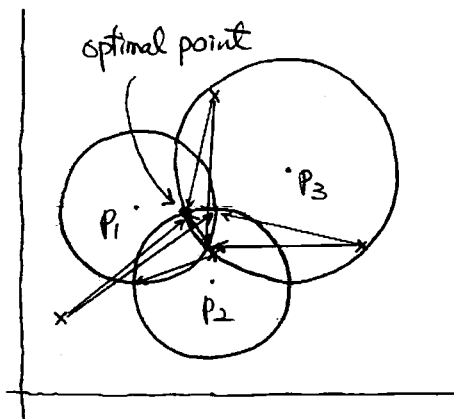
cell $R_3(1, -1, 1)$ $x = 5/1 = 5$, $y = 8/1 = 8$ $(x, y) = (5, 8)$

cell $R_4(-1, 1, -1)$ $x = (-5)/(-1) = 5$, $y = (-8)/(-1) = 8$ $(x, y) = (5, 8)$

cell $R_5(-1, -1, -1)$ $x = (-5)/(-3) = 5/3$, $y = (-14)/(-3) = 14/3$ $(x, y) = (5/3, 14/3)$

cell $R_6(1, -1, -1)$ $x = -9/(-1) = 9$, $y = -4/(-1) = 4$ $(x, y) = (9, 4)$

cell $R_7(1, 1, -1)$ $x = 1/1 = 1$, $y = 2/1 = 2$ $(x, y) = (1, 2)$



No apex point lies in the interior of the corresponding cell.

for each cell, find a nearest/farthest point from the apex point

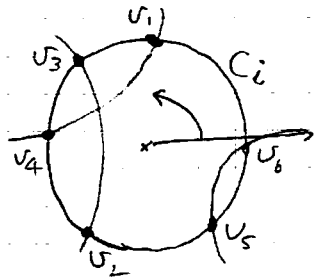
Construct an arrangement of circles.

① Compute all intersections of circles

→ $O(n^2)$ intersections

we compute a list of all intersections for each circle

circle C_i $\boxed{v_1 | v_2 | \dots | v_k}$ ← list of intersections on C_i



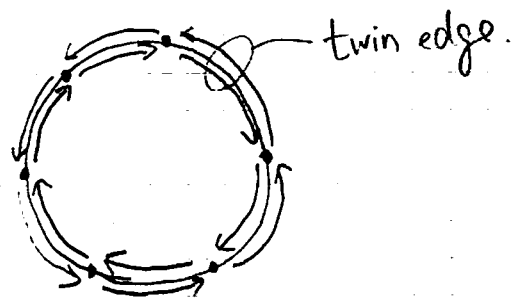
Sort all the intersections by angles from the horizontal ray to the right from the center of C_i .

$(v_{i_0}, v_{i_1}, \dots, v_{i_{k-1}})$: sorted list

→ make edges $e_{i_0} = (v_{i_0}, v_{i_1}), \dots, e_{i_{k-1}} = (v_{i_{k-2}}, v_{i_{k-1}})$,

$e_{i_k} = (v_{i_{k-1}}, v_{i_0})$

→ for each edge, we create its twin edge of the opposite direction



Now we have all vertices (intersections) and edges connecting vertices.

Next, we enumerate all the cells

while (there is an unvisited edges) {

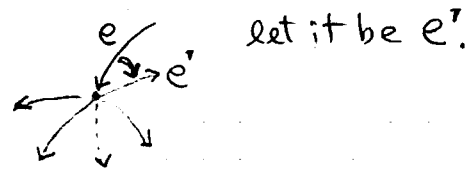
 Choose one unvisited edge e arbitrarily;

 Make a new cell R_i lying to the left of e ;

 Put the edge into a list $L(R_i)$ associated with R_i .

 do {

 Among those edges outgoing from the target point of the edge e , choose the one that is next to e in the clockwise order.



 Put e' into the list $L(R_i)$ and label it as "visited".

$e = e'$;

 while (the list $L(R_i)$ is not a cycle);

}

Algorithm

- (1) Draw n circles C_1, \dots, C_n ;
- (2) Construct an arrangement of those circles.
- (3) Compute a dual graph G ;
- (4) Compute a spanning tree T of G ;
- (5) Compute a walk along T ; \rightarrow sequence of cells.

(6) R = initial cell in the sequence

Compute the objective function S

(compute coefficients A, B, C , and D

$$A(x^2 + y^2) + Bx + Cy + D).$$

Solve a system of equation $\partial S / \partial x = \partial S / \partial y = 0$;

let $p(R)$ be a point satisfying the system;

Check whether $p(R)$ lies inside R or not
by traversing the boundary of R ;

if $p(R)$ lies inside R then $p^*(R) = p(R)$;

else if $A > 0$ then

find a point on the boundary of R that
is closest to $p(R)$

else $A < 0$ then

find a point on the boundary of R that
is farthest from $p(R)$

else $A = 0$ then

find a point on the boundary of R that
minimizes $Bx + Cy + D$;

all by traversing the boundary of R ;

Now, we have a dual graph $G = (V, E)$.

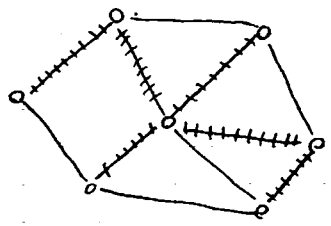
of vertices = # of cells = $O(n^2)$

of edges \leq # of twin edges = $O(n^2)$

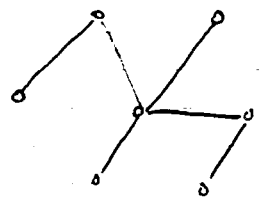
time to compute the dual graph = $O(n^2)$.



Then, compute a spanning tree T of G



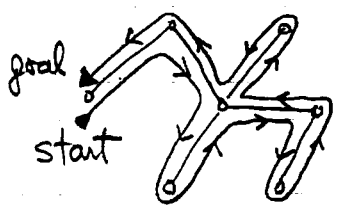
dual graph



tree edge.



Define a walk along T



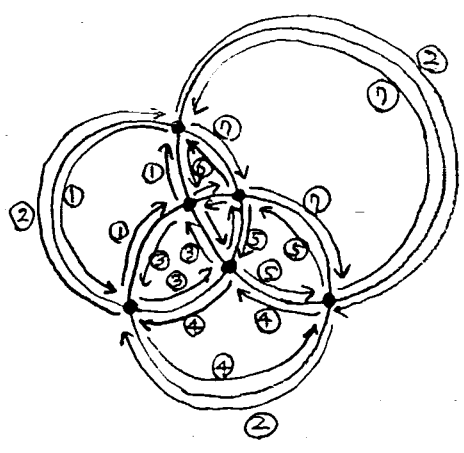
Walk

= a sequence of vertices visiting each edge exactly twice.

two vertices are adjacent in a walk

\Rightarrow their corresponding cells are adjacent

(Example E9)



Now, we can traverse the boundary of any cell.

Constructing a dual graph

Dual graph $G = (V, E)$

vertex \leftarrow cell.

edge \leftarrow between two cells adjacent to each other.

For each cell R_i

traverse the boundary of R_i

for each edge e

find its twin edge e'

let R_j be the cell associated with e'

define an edge (R_i, R_j) if it is not found.

while (R is not the last cell in the sequence) {
 let R' be the next cell in the sequence;
 modify the objective function using
 the information of the edge between R and R';
 (this modification can be done in $O(1)$ time)
 let $R = R'$;
 repeat the same calculation as above;
 → find an optimal point $p^*(R)$ for R;
 }

Note that the boundary of a cell is traversed only twice. Thus, the total time for traverse is $O(n^2)$.

Therefore, the total running time is
 $O(n^2 \log n)$ for sorting intersections
 $O(n^2)$ for remaining computation
 in total $O(n^2 \log n)$ time and
 $O(n^2)$ space



can improve the running time into $O(n^2)$?

space can be improved $O(n)$
 while keeping the running time in $O(n^2 \log n)$
 ← Plane Sweep Algorithm.